



HOME SITE DELIVERY REQUEST FORM

STORE: _____

P.O. #: _____

SALESPERSON: _____

ORDER DATE: _____

CUSTOMER INFORMATION (please print)

NAME: _____

DAY PHONE #: _____

DELIVERY ADDRESS: _____

EVENING PH. #: _____

*PLEASE CHECK ONE ONLY

Customer Has Ordered Cabinets
(Due on _____.)

No Cabinets.

CUSTOMER ACKNOWLEDGEMENTS & DELIVERY RESTRICTIONS

1. Appointments for delivery will be made once the order is complete and must be accepted within (2) weeks of completion.
2. Deliveries will be made Monday through Friday 8am - 12pm or 12pm - 4pm on the day the truck is in the delivery area. Note: Customer must be flexible with our delivery schedule on time of day and day of week.
3. Delivery is limited to curbside, porch, or garage placement only. **NO SECOND FLOOR OR BASEMENT DELIVERIES.**
4. Driver reserves the right to decline delivery to a location which may result in damage to merchandise, property or cause bodily harm. We will not put tops through windows, bulkheads, up ladders, etc.
5. Customer is responsible for accepting and signing for countertops in good condition. (No claims for damages will be accepted after delivery.)
6. In the event that the customer is not at the delivery address at the appointment time, the countertops will be returned to our warehouse and are subject to a store delivery or second delivery fee.
7. Counter Pro, Inc.® is not responsible for delays in shipment of merchandise. No claims for lost wages or subcontractor fees incurred will be accepted or reimbursed.

I agree to the above delivery conditions. (*Signature must be customer's only.*)

Customer Signature: _____ Date: _____

*****THIS FORM MUST ACCOMPANY THE COUNTERTOP ORDER, OR IT WILL BE SHIPPED DIRECTLY TO THE STORE.

STANDARD MODEL IN LOGIC

THE COMPLETION THEOREM FOR FIRST ORDER LOGIC

Let \mathcal{L} be a first order language. A set of sentences Σ in \mathcal{L} is called **consistent** if there is no sentence ϕ such that $\Sigma \cup \{\phi\}$ is unsatisfiable. A set Σ is **maximally consistent** if it is consistent and for every sentence ϕ , either $\phi \in \Sigma$ or $\neg\phi \in \Sigma$.

Let Σ be a consistent set of sentences. A **model** for Σ is a structure \mathcal{M} such that $\mathcal{M} \models \Sigma$. The **Skolemization** of Σ is the set Σ^s obtained by replacing each existential quantifier $\exists x$ by a Skolem function symbol f .

Let \mathcal{M} be a model for Σ . A **Skolem hull** H of \mathcal{M} is a substructure of \mathcal{M} such that $H \models \Sigma^s$ and H is closed under the Skolem functions.

Let \mathcal{M} be a model for Σ . A **Skolem model** for Σ is a Skolem hull H of \mathcal{M} .

Let Σ be a consistent set of sentences. Then there exists a Skolem model for Σ .

THE COMPLETION THEOREM FOR FIRST ORDER LOGIC

Let Σ be a consistent set of sentences. Then there exists a model for Σ .

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