

Trust Without Disclosure: From the Pub to the Protocol

Why Zero-Knowledge Proofs are the Privacy Backbone for DeFi on Blockchain

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Over 18? Prove it! (but keep the rest of your data private)

Imagine you're asked for ID at a pub.

To prove you're over 18, you show your driver's license, and with it, you expose your **full name**, **address**, **birthday and more**.

All of that, just to confirm one binary truth:

"Yes, I'm old enough."

This is the kind of data oversharing we replicate every day in digital systems, especially on blockchain.

What if you could prove your age, without revealing it?

That's what **Zero-Knowledge Proofs (ZKPs)** enable.

Now, imagine you're applying for a **loan on the blockchain** (in a world of DeFi, decentralised peer to peer lending), or **selling carbon credits** (in a world of decentralised carbon trading).

To meet lending or trading criteria, you may need to prove things like:

- Vour wallet has over \$5M and you have secured sufficient collateral
- Vou have a high enough credit score
- Vou have passed KYC compliance checks
- Vour carbon credit is authentic

Should you have to reveal every transaction you've ever made to your counterparty, or in the case of public permissionless blockchains, the whole world? Should you have to reveal all personal details?

That's where **Zero-Knowledge Proofs** (**ZKPs**) come in.

Note: This article does not address what roles financial and related institutions could play in that decentralised world (e.g. authentication, collateral, escrow or advisory services etc). This is best served in dedicated article for blockchain-based decentralised lending. However, it does deal with how only the minimum information could be shared on chain, leveraging smart contracts, without requiring a centralised intermediary.



What is a Zero-Knowledge Proof?

A ZKP allows someone to prove a statement is true (like age > 18) without revealing the underlying data (your actual age).

This isn't magic. It's maths, and it's how we go from:

- X "Share everything to prove anything"
- V to "Prove only what's necessary, and nothing more"

A ZKP lets you prove something is true, without revealing without revealing underlying data.

- **V** "I am over 18" (without showing your birthdate)
- V"I have over \$5M" (without showing your wallet balance)
- V "I have a high enough credit score" (without showing your credit score or history)
- V "I have passed KYC compliance checks" (without revealing your personal details)
- **V** "This carbon credit is real" (without revealing your supplier)

Mathematics and cryptography enable a new kind of trust for decentralised, blockchain based systems.

From the Pub: Simulating Zero-Knowledge in Python

Let's recreate that "over 18" scenario in code (simplified hashing approach, just to prove the concept):

```
from hashlib import sha256
# Prover's secret (not to be made visible to the Verifier)
your age = 34
threshold = 18
# Commitment (a hash of the age, to be made visible to the Verifier)
def create commitment(value):
    return sha256(str(value).encode()).hexdigest()
commitment = create commitment(your age) # visible to the Verifier
# Prover generates a proof (input age not made visible to the Verifier, only
output and commitment)
def generate proof(age, threshold, commitment):
    if age > threshold and create commitment(age) == commitment:
        return {"proof_passed": True, "commitment": commitment}
    else:
       return {"proof passed": False, "commitment": None}
proof package = generate proof(your age, threshold, commitment) # output
visible to the Verifier
# Verifier sees only proof, not age
def verify_proof(proof):
    return proof["proof passed"] and proof["commitment"] == commitment
if verify proof (proof package):
    print("✓ Verified: Over 18 proven without revealing age.")
   print("X Verification failed.")
```

The verifier **never sees the age**, just the commitment and proof result (and the logic executed).

"Give me a hash of your age. Now prove to me that whatever number you hashed is over 18, without telling me what the number is."

A Zero-Knowledge Proof allows the **prover** to commit to a secret and generate a **verifiable proof**, without ever revealing the underlying data.

11 To the Protocol: ZKPs in Decentralised Lending

Let's now move from the pub to a **DeFi protocol**.

Suppose a protocol requires that you:

- ✓ Hold more than \$5M
- W Have no history of default
- **Qualify** for credit

Let's assume that sharing your wallet address, full transaction history, or asset mix is a **nonstarter** for institutions, DAOs, and privacy-conscious individuals.

First, let's move from Python Hashing to Elliptic Curve Pedersen Commitments

In the earlier example, we used Python's built-in sha256 to simulate a commitment:

```
def create_commitment(value):
    return sha256(str(value).encode()).hexdigest()
```

This gives the idea of a commitment (i.e. hiding a secret value while being "locked in") but it falls short for real-world zero-knowledge applications.

■ Limitations of SHA256 (Hash-Based Commitments)

Weakness	Why It Fails in Real ZK Contexts
X No randomness	Without blinding, it's guessable (e.g. small ages or balances)
X Not zero- knowledge	Hashes reveal structure (e.g. collisions, entropy leakage)
X Not algebraic	You can't prove logic (like x > 5M) over hash values in Zero-Knowledge Succinct Non-Interactive Argument of Knowledge (zkSNARK) circuits
X Not compatible	zkSNARKs work over elliptic curves, not hashes
X Not composable	You can't homomorphically add or manipulate SHA256-based commitments

☑ Enter Elliptic Curves and Pedersen Commitments

In modern ZKPs, we use algebraic commitments. The most common is the Pedersen commitment, which is built on Elliptic Curve Cryptography (ECC):

$$C = g^x \cdot h^r \mod p$$

This is the multiplicative form, based on integer finite field maths. It can be used in **non-Elliptic** Curve settings, where g and h are integers (generators in a finite multiplicative group), x = secret value; r = random blinding. This is more easily simulated in Python, using pow() and % p, for educational and quick simulation purposes. It is not used in real ZK circuits.

In **Elliptic Curve** terms, this is expressed as:

$$C = x \cdot G + r \cdot H$$

Where:

- x = the secret (e.g. wallet balance)
- r = random blinding factor
- G, H = base points (generators) on an elliptic curve
- C = the resulting **elliptic curve point** (i.e. the commitment)

☑ Why Elliptic Curves?

Elliptic curves are the foundation of modern cryptography and Zero-Knowledge Succinct Non-Interactive Argument of Knowledge (zk-SNARK) because they offer:

- **Strong security** with short keys
- **Fast arithmetic** over finite fields
- **+ Group structure**, which lets us "add" and "multiply" secrets in a proof system
- Compatibility with zero-knowledge systems like Groth16, PLONK, Halo2, and Bulletproofs

In ZKPs, these curve-based commitments:

- Hide the data completely (even from the verifier)
- Let the prover **demonstrate facts** about the data (like "I have > \$5M") without ever revealing it
- Work inside arithmetic constraint systems used in zkSNARK circuits

In Short:

- **Python hash** = good for learning, bad for privacy and ZKP
- **Pedersen on elliptic curves** = private, provable, zkSNARK-compatible

It's the difference between sealing your secret with tape versus locking it in a zero-knowledge vault, built from elliptic curve cryptography.

Pedersen commitment-based proof using elliptic curve cryptography, using bulletproofs crate in Rust

The code below is realistic implementation of a Pedersen commitment-based solvency proof, using **bulletproofs** crate, built on elliptic curve cryptography (Curve25519 via the Ristretto group) in Rust.

bulletproofs crate is a Rust cryptographic library that allows developers to create zero-knowledge proofs, specifically range proofs and constraint systems, without revealing the underlying data. In the code, it enables us to prove that a private balance is greater than a public threshold, without disclosing the actual balance.

It uses Pedersen commitments to securely hide values, and Rank-1 Constraint Systems (R1CS) to define mathematical rules those values must satisfy. R1CS lets you turn a program (or logic) into a series of equations that must be true.

Crucially, **bulletproofs** require no trusted setup and produce small, efficient proofs ideal for privacy-preserving blockchain applications

The protocol ensures the verifier only sees a **proof of solvency**, but never the wallet balance.

This is the foundation of **zk-DeFi**: trusted, verifiable, and private-by-design.

toml (Cargo.toml)

```
[package]
name = "zk_solvent_proof"
version = "0.1.0"
edition = "2021"

[dependencies]
bulletproofs = "4.0.0"
curve25519-dalek = "4.1.1"
merlin = "3.0.0"
rand = "0.8"
```

```
// Proves in zero-knowledge that someone's balance is greater than a public
threshold
// WITHOUT revealing the actual balance - using Bulletproofs (no trusted setup).
         bulletproofs::r1cs::{Prover,
                                          Verifier,
                                                        ConstraintSystem,
LinearCombination, R1CSProof, Variable};
use bulletproofs::{BulletproofGens, PedersenGens}; // Generator systems for
commitments and proofs
     curve25519 dalek::ristretto::CompressedRistretto; // For outputting
commitments
use curve25519 dalek::scalar::Scalar; // Finite field numbers
use merlin::Transcript; // Keeps context consistent between prover and verifier
use rand::rngs::OsRng; // Secure randomness
/// This function is used by the prover (the person who wants a loan, for
example).
/// It generates a zero-knowledge proof that they have more money than a certain
/// But without showing exactly how much money they have.
fn generate solvent proof(
   threshold: u64
                     // Public threshold (e.g., minimum required to borrow)
) -> (R1CSProof, CompressedRistretto, PedersenGens, BulletproofGens) {
   // Standard generator setups - these define the cryptographic playground
   let pc gens = PedersenGens::default();
                                                // For commitments (hiding
numbers)
   let bp gens = BulletproofGens::new(64, 1);  // For proving in 64-bit
range
   let mut rng = OsRng; // Secure random number generator
   // Compute the difference between balance and threshold
   let delta = balance - threshold;
   assert! (delta > 0, "Balance must exceed threshold"); // This is what we're
proving
   // Create a transcript to ensure both parties are on the same page
cryptographically
   let mut transcript = Transcript::new(b"ZK Solvency Proof");
   let mut prover = Prover::new(&pc gens, &mut transcript);
   // === STEP 1: COMMIT TO BALANCE AND THRESHOLD ===
   // These commitments hide the values but let us prove things about them.
   let (com balance, var balance) = prover.commit(Scalar::from(balance),
Scalar::random(&mut rng));
         ( , var threshold) = prover.commit(Scalar::from(threshold),
Scalar::zero()); // threshold is public, so no blinding
```

```
// === STEP 2: COMMIT TO DELTA ===
           (com delta,
                        var delta) = prover.commit(Scalar::from(delta),
Scalar::random(&mut rng));
    // === STEP 3: ENFORCE THE RULE ===
    // balance = threshold + delta
    prover.constrain(var_balance - var_threshold - var_delta);
    // === STEP 4: PROVE THAT delta > 0 ===
    // We do this by showing delta is in a valid range [1, 2^32]
    let mut exp = Scalar::one(); // Tracks 2^i
    let mut delta lc = LinearCombination::default(); // Linear sum of bits *
powers of 2
    for i in 0..32 {
        let bit = (delta >> i) & 1;
       let ( , var bit) = prover.commit(Scalar::from(bit), Scalar::random(&mut
rng));
       prover.constrain(var bit - var bit * var bit); // Force the bit to be
0 or 1
       delta lc = delta lc + (exp, var bit); // Add to the delta representation
       exp = exp + exp;
    }
    // Make sure all bits sum back to delta
    prover.constrain(var delta - delta lc);
    // === STEP 5: GENERATE THE PROOF ===
    let proof = prover.prove(&bp gens).expect("Proof generation failed");
    // Return: proof, commitment to balance, and the gens so verifier can use
them
    (proof, com balance.compress(), pc gens, bp gens)
/// Verifier logic: checks the proof, using only the proof and a commitment.
/// It doesn't see the balance, but can verify the relationship is true.
fn verify solvent proof(
    proof: R1CSProof,
    com balance: CompressedRistretto,
    pc gens: PedersenGens,
   bp gens: BulletproofGens,
) -> bool {
    let mut verifier transcript = Transcript::new(b"ZK Solvency Proof");
    let mut verifier = Verifier::new(&mut verifier transcript);
    // Placeholder for commitment to balance
    let var balance = verifier.commit(com balance);
```

```
// The verifier doesn't know the threshold or delta, so it uses dummy zero
commitments
   let
         var threshold =
                               verifier.commit(pc gens.commit(Scalar::zero(),
Scalar::zero()).compress());
          var delta
                               verifier.commit(pc gens.commit(Scalar::zero(),
Scalar::zero()).compress());
   // Enforce same logic: balance = threshold + delta
   verifier.constrain(var balance - var threshold - var delta);
   // Enforce that delta > 0 via same 32-bit decomposition (dummy bits)
   let mut exp = Scalar::one();
    let mut delta lc = LinearCombination::default();
    for in 0..32 {
        let
             var bit
                        = verifier.commit(pc gens.commit(Scalar::zero(),
Scalar::zero()).compress());
       verifier.constrain(var bit - var bit * var bit); // Bit must be 0 or 1
       delta lc = delta lc + (exp, var bit);
       exp = exp + exp;
   verifier.constrain(var delta - delta lc);
   // If everything checks out, the proof is valid
   proof.verify(&bp gens, &pc gens, verifier).is ok()
}
fn main() {
    // Example use case: someone has 6.2 million, threshold is 5 million
    let balance = 6 200 000u64;
    let threshold = 5 000 000u64;
    // Prover generates a proof
    let.
              (proof,
                          com balance,
                                              pc gens, bp gens)
generate solvent proof(balance, threshold);
    // Verifier checks the proof - without seeing the actual balance
    let valid = verify solvent proof(proof, com balance, pc gens, bp gens);
    if valid {
       println!("♥ ZK Proof verified successfully.");
    } else {
       println!("X ZK Proof verification failed.");
   // Display the cryptographic commitment to the balance
   println!(" Balance commitment: {:?}", com balance);
}
```

This Rust code proves something very powerful:

- The prover (say, a DeFi borrower) can prove they have more than a minimum required balance (e.g., for a loan)
- They do not reveal their actual wallet balance
- The verifier (say, a smart contract or DeFi protocol) can cryptographically verify the
- This is done with **no trusted setup** and **small proof size** using **Bulletproofs**



Nhy ZKPs Matter for DeFi Infrastructure

Without ZKPs:

- Transparency becomes exposure.
- Compliance becomes surveillance.
- Trust becomes a liability.

With ZKPs:

- **Users can prove eligibility without doxxing wallets.**
- **All Protocols can enforce rules without reading your data.**
- **1** Institutions can interact on-chain without leaking IP or identity.

ZKPs are the **privacy layer DeFi needs**, not to hide, but to scale.



What I've Seen in the Field

ZKPs can be used to:

- Prove carbon credit validity without revealing project-level identities
- V Enable confidential ESG financing
- V Build pilots where Ethereum smart contracts verify zk-SNARK proofs using Solidity and Rust

ZKPs are already making decentralised systems compliant, credible, and confidential.

Zero-Knowledge Proofs in International Standards

- ISO/IEC 27565 provides internationally recognized guidelines on the use of zero-knowledge proofs (ZKPs) to enhance privacy during data sharing. This work is part of the broader effort led by ISO/IEC JTC 1/SC 27, the subcommittee responsible for global standards in cybersecurity, cryptography, and privacy protection.
- Meanwhile, ISO TC 307, which focuses on blockchain and distributed ledger technologies (DLT), recognizes privacy and security as central challenges in decentralized systems. Within TC 307, Working Group 2 (WG2) and the Joint Working Group 4 (JWG 4) actively explore privacy-preserving mechanisms, including zero-knowledge proofs.
- Notably, technical reports such as ISO/TR 23244 and ISO/TR 23455 outline principles for privacy and smart contract security that naturally align with ZKP-based approaches.

Further Reading & Technical Sources

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